

5. A NON-EXISTENT HIERARCHY

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complete when so operated on. A most convenient starting point is the following family T which is of p -bounded formula size:

$$T_{n^2+n} = \prod_{k=1}^n \sum_{i=1}^n x_{k,i} y_i .$$

Clearly (i) the coefficient of $y_1 \dots y_n$ in T_{n^2+n} ,

$$(ii) \quad \frac{\partial}{\partial y_1} \frac{\partial}{\partial y_2} \dots \frac{\partial}{\partial y_n} T_{n^2+n} , \text{ and}$$

$$(iii) \quad \left(\frac{3}{2}\right)^n \int_{-1}^1 \dots \int_{-1}^1 [y_1 \dots y_n T_{n^2+n}] dy_1 \dots dy_n$$

all equal $\text{Perm} \{ x_{k,i} \}$.

In contrast, it is easy to see that all the other operations that we have considered preserve p -computability. This is immediate in the case of substitution. It can be shown to be true for $\partial P / \partial x_i$ and $\int P dx_i$ by considering a program for P , and decomposing it according to the powers of x_i at each instruction in the manner of [12].

5. A NON-EXISTENT HIERARCHY

By analogies with recursion theory we can attempt to define the following hierarchy:

Definition. PD^0 = class of p -computable polynomial families. For $i > 0$ $P \in PD^i$ iff P is defined by some $Q \in PD^{i-1}$ in the sense of Definition 3.

That this hierarchy collapses in this algebraic case is easy to see:

THEOREM 5. For any F and any $i > 0$ $PD^i = PD^{i+1}$.

Proof. It is clearly sufficient to prove $PD^1 = PD^2$. If $P \in PD^2$ then for each m

$$P_m(\mathbf{x}) = \sum_{\mathbf{b}} Q_i(\mathbf{x}, \mathbf{b})$$

where for some $R \in PD^0$ for each i

$$Q_i(\mathbf{x}, \mathbf{b}) = \sum_{\mathbf{c}} R_j(\mathbf{x}, \mathbf{b}, \mathbf{c}) .$$

Hence

$$P_m(x) = \sum_{\mathbf{b}, \mathbf{c}} R_j(\mathbf{x}, \mathbf{b}, \mathbf{c})$$

which shows that $P \in PD^1$. □

We can attempt to generalise the definition of the above vacuous hierarchy by allowing the number of "alternations" to increase with the number of indeterminates.

Let t be any polynomial. Define $t-D^0$ to be the class of t -computable families. For $i > 0$ let $t-D^i$ be the class of families that are defined by some family in $t-D^{i-1}$ in the sense of Definition 3. Finally PD^* is the class of all families P such that for some t

$$P = \{P_i \mid P_i = Q_i \text{ for some } Q \in t-D^{(i)}\}.$$

THEOREM 6. $PD^* = PD^1$

Proof. Similar to previous theorem. □

The above two results should be contrasted with the Boolean case where they still hold formally, but are no longer natural. The above definition of the successive levels PD^i is only natural if each level is a robust closure class. In Boolean algebra, however, PD^i is not known to be closed under complementation for any $i \geq 1$. Analogues of PD^i and PD^* where complementation is allowed at each level of alternation are not known to collapse, and are merely finite versions of the Meyer-Stockmeyer hierarchy, and PSPACE respectively [10].

A simple application of Theorem 5 is in recognising such polynomials as $\#HG$ as being p -definable. An intriguing open question is whether HG itself is p -definable for each F . If it is not then $P \neq NP$ (see Proposition 4 in [13]). If it is then the Meyer-Stockmeyer hierarchy and PSPACE can be simulated within p -definable families of polynomials.

6. UNIVERSALITY OF LINEAR PROGRAMMING

Here we consider a Boolean function family LP that corresponds to a linear programming problem and show that every p -computable family is the p -projection of it. Thus for computing discrete functions in polynomial time a package for LP for each input size is sufficient and no further programming is required. If we fix certain of the arguments of LP_i according to the particular function and input size being computed, the package becomes a program for the required function. That LP is itself p -computable follows from the recent result of Khachian [8].

The reader should note that several tractable problems in combinatorial optimisation are already known to have linear programming formula-